



Recent Algorithmic Developments in NetworKit

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Agenda



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- numerous algorithmic additions to NETWORKIT
- and lots of refactoring (see previous talk)



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In this talk: new algorithms and features since ND'17

(Brief tour through various modules ...)



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Use cases:

- Null-model for network analytics (e.g., modularity)
- Benchmarking graph algorithms





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In NETWORKIT: Curveball and GlobalCurveball (pick many trades at a time)





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- Top-k (harmonic) closeness [BBCMM16] Computes k vertices with highest closeness w/o computing all scores





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In NETWORKIT: fastest available betweenness approximation





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- Computation of group centrality scores
- Finding groups with maximal centrality (usually a hard problem)





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 Upcoming: Group closeness [AvdGM19] Fast local-search algorithm





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- Varint encoding: bit-length of IDs adapted to # nodes of the graph
- Goal: represent all data available in NETWORKIT(weights, IDs, ...) in a compact format



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a.k.a. Representation Learning

Problem: Given graph *G*, map each vertex $v \in V(G)$ to $f(v) \in \mathbb{R}^d$ for some *d*.



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Other embedding algorithms in the future?



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 Classical graph problem; useful building block for other algorithms



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- Graph generator by F.-B. Mocnik [M18] Models spacial graphs













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Available from: https://github.com/hu-macsy/simexpal





Thank You!

